

Consistency between transitive relations and between cones

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Abstract

A relation extends another relation consistently if its symmetric, respectively its asymmetric part contains the corresponding part of the smaller relation. It is shown that there exists no finite circular chain made from two transitive relations \mathbf{A} and \mathbf{B} with at least one link from their asymmetric parts if and only if there exists a total preorder which consistently extends both. This extension is uniquely determined if and only if the reflexive transitive closure of the union of \mathbf{A} and \mathbf{B} is total. Applications: (1) If the steps of a walk come from two positive cones, with at least one step from one of the cones' non-linear parts, then returning to the origin is impossible if and only if there exists a third cone of which the linear part contains each of the linear parts of the two original cones, and of which the non-linear part contains each of the two non-linear parts. (2) Reminiscent of the Fundamental Theorem of Asset Pricing, absence of arbitrage is equivalent to the existence of a complete preference order which consistently extends a market's fair exchange relation and its objective strict preference order. (3) Another impossibility in microeconomics: For two agents with additive positively homogeneous preferences, an allocation can never be Pareto optimal if they are not 'cut from the same cloth'.

Key words: Cone, preference order, preorder, strict partial order, transitive relation, vector order.

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1 Intentions

Any relation splits into its symmetric and its asymmetric part. For instance, a preorder is the disjoint union of an equivalence relation and a strict partial order. While this is trivial,

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it is easy to find examples of disjoint equivalence relations and strict partial orders whose union is neither a preorder, nor even just a transitive relation. The question of what could manifest a meaningful reverse direction of the splitting procedure for – more generally than preorders – transitive relations therefore seems to deserve closer examination.

In my literature research, I could not find any related previous results. Apart from partial orders, transitive relations or preorders generally do not seem to draw a lot of attention in mathematics. At least in economic theory, preorders do, and are usually referred to as preference orders. However, the focus there typically lies on questions that I will not consider here (e.g. [6]). Textbooks – as extensive and thorough they may be on orders and lattices (e.g. [8] or [11]) – seem to mostly skip over the topic of transitive relations and preorders, as well.

Diving further into the subject, it became clear to me that to meaningfully construct a preorder from a given equivalence relation and a strict partial order first of all means that the given equivalence relation needs to be contained in (i.e. imply) the one derived from (or implied by) the constructed preorder, and, similarly, the given strict partial order needs to be contained in (imply) the derived one. Another observation was that, for the given equivalence relation and the given strict partial order, no finite closed (i.e. circular) chain containing both relations should exist, as this in a quite obvious sense means that they are not consistent with one another. It turns out, the two conditions are equivalent.

More generally, my main theorem states that two transitive relations are chain-consistent with one another in the sense that there exists no finite circular chain made from them with at least one link from their asymmetric parts if and only if there exists a total preorder which consistently extends both of them in the sense that the symmetric, respectively the asymmetric part of each of the two given relations lies in the corresponding part of the extending preorder. Expressed differently, two transitive relations are consistent if and only if they both can be meaningfully embedded into a third one which also possesses a completeness property. Additionally, the extending third relation is uniquely determined if and only if the reflexive transitive closure of the union of the two original relations is total.

My methods stay elementary, because relations are very fundamental, but they also are very basic mathematical objects. As with some other extension results, such as Hahn-Banach, no deep theorems are required. Still, my results contribute to a deeper understanding of transitive relations, and – more specifically – to the understanding of the relationship between preorders, equivalence relations, and strict partial orders.

I then provide three applications of my theorem. One in geometry or topological vector spaces, one in financial economics, and the third in microeconomics.

The first one is related to the fact that a walk with steps taken from a positive cone cannot end at the starting point if at least one step is from the cone's non-linear part, which is the cone less the largest linear space contained in it. What if the steps come from two cones? My main theorem implies that if the steps of a walk come from two positive cones, with at least one step from one of the cones' non-linear parts, then returning to the origin is impossible if and only if there exists a third cone of which the linear part contains each of the linear parts of the two original cones, and of which the non-linear part contains each of the two non-linear parts of the original cones. I also formulate this as follows: Two positive cones are path-consistent if and only if there exists a common

consistent extension into a third cone. In an example, I then illustrate how a related notion of (path-)consistency between a cone and a linear functional implies a strict type of positivity of that functional with respect to the cone.

The second example concerns financial economics or mathematical finance. I show that if a market is modeled by an equivalence relation for its objectively fair exchanges, meaning an indifference relation, and by a strict partial order for its objective down-trades, meaning an objective strict preference order, then such a market is free of arbitrage opportunities, or so-called free lunches, if and only if there exists a complete preference order which consistently extends the fair exchange and the down-trade relation. I make the argument that this result is so reminiscent of the seminal Fundamental Theorem of Asset Pricing, which also is an extension theorem, that it can be considered a broad generalization of it.

The third application is in microeconomics: For two agents, or an agent and a market, with additive positively homogeneous preferences, meaning vector preorders, an allocation can never be Pareto optimal if the agents' preferences are not consistent with one another in the earlier explained sense. Expressed differently, if both agents' preferences cannot meaningfully, i.e. consistently, be embedded into the same bigger preference order, they will never reach an equilibrium.

While very straightforward, I could not find most of the needed preliminary results on transitive relations or preorders in the required generality in the literature, so I included them with proofs. Where possible, I refer to the literature even if the reference only points to a related result for partial orders.

2 Basic properties of relations

For the sake of notation and for the convenience of the reader, in this section, I recapitulate several basic definitions and properties of relations, as well as several basic results on them. In particular, the generation of relations with certain properties is an important and later recurring topic. In the following, all relations on a non-empty set \mathcal{M} are understood to be homogeneous binary relations, meaning they are non-empty subsets of $\mathcal{M} \times \mathcal{M}$.

DEFINITION 1. *A relation \mathbf{R} on a set \mathcal{M} can have one or several of the following properties. For any $\alpha, \beta, \gamma, \delta \in \mathcal{M}$ and for $\alpha \mathbf{R} \beta$ denoting $(\alpha, \beta) \in \mathbf{R}$:*

1. *Transitivity: If $\alpha \mathbf{R} \beta$ and $\beta \mathbf{R} \gamma$, then $\alpha \mathbf{R} \gamma$.*
2. *Symmetry: $\alpha \mathbf{R} \beta$ implies $\beta \mathbf{R} \alpha$.*
3. *Asymmetry: $\alpha \mathbf{R} \beta$ implies $\beta \not\mathbf{R} \alpha$.*
4. *Reflexivity: $\alpha \mathbf{R} \alpha$.*
5. *Irreflexivity: $\alpha \not\mathbf{R} \alpha$.*
6. *Antisymmetry: $\alpha \mathbf{R} \beta$ and $\beta \mathbf{R} \alpha$ implies $\alpha = \beta$.*
7. *Additivity: For \mathcal{M} a real linear space, $\alpha \mathbf{R} \beta$ and $\gamma \mathbf{R} \delta$ implies $(\alpha + \gamma) \mathbf{R} (\beta + \delta)$.*

8. *Positive homogeneity:* For \mathcal{M} a real linear space, $\alpha \mathbf{R} \beta$ and $r \in \mathbb{R}_0^+$ implies $(r\alpha) \mathbf{R} (r\beta)$.

9. *Totality:* Either $\alpha \mathbf{R} \beta$, $\beta \mathbf{R} \alpha$, or both.

Obviously, properties of Def. 1 can be mutually exclusive, or a combination of properties can imply another one. For instance, reflexivity with additivity implies transitivity. At least in the specific case of partial orders (properties 1, 4, and 6; e.g. [11]), the proposition below is textbook knowledge.

PROPOSITION 1. *Let \mathcal{R} be a family of relations on \mathcal{M} which all share the same non-empty subset of the first eight properties of Def. 1. Assume that \mathcal{M} is a real linear space if properties 7 or 8 are considered. If $\bigcap \mathcal{R}$ is non-empty, then it is a relation on \mathcal{M} with the same set of properties that the individual relations have.*

Proof. For any single one of the properties 1 to 8, if all relations have the corresponding property, then it is straightforward to check that a non-empty intersection must have it. \square

COROLLARY 1. *Let \mathbf{S} be an arbitrary non-empty subset of $\mathcal{M} \times \mathcal{M}$ and choose a non-empty subset of the properties 1, 2, 4, 7, and 8 of Def. 1. Assume that \mathcal{M} is a real linear space in case properties 7 or 8 are chosen. Then there exists a uniquely determined smallest relation on \mathcal{M} that contains \mathbf{S} and that has the chosen properties.*

Smallest here means an inclusion-minimal relation.

Proof. Since the trivial relation, $\mathcal{M} \times \mathcal{M}$, has the properties 1, 2, 4, 7, and 8, the intersection of all relations on \mathcal{M} which contain \mathbf{S} and have the chosen properties is non-empty. By Prop. 1, this intersection again has the chosen properties. By construction, there is no smaller such relation. \square

DEFINITION 2. *The uniquely determined smallest relation of Corollary 1 is called the relation with the chosen properties on \mathcal{M} generated by \mathbf{S} .*

From now on, transitive relations will be denoted by the symbol \preceq , which can come with various indices, or other attached symbols, to create distinctive notation.

DEFINITION 3. *A transitive relation \preceq is a relation with property 1 of Def. 1. For an arbitrary non-empty $\mathbf{S} \subset \mathcal{M} \times \mathcal{M}$, let $\preceq_{\mathbf{S}}$ denote the transitive relation generated by \mathbf{S} . If \mathcal{M} is a real linear space, let $\preceq_{\mathbf{S}}^{add,ph}$ denote the transitive additive positively homogeneous relation generated by \mathbf{S} .*

There is a reason why I positioned transitivity at the top in Def. 1. If I think of any relation that is supposed to have some ordering capability, transitivity seems to be the most basic property I would expect: If some β is at least as great or as good as some α , and some γ is at least as good or as great as β , then I should reasonably be able to expect γ to be at least as great or good as α . With regard to (strictly) ordering anything, there seems to be little value to a relation that is not at least transitive. This attitude is, of course, reflected by the presence of transitivity in the contemporary definitions of preorders, partial orders, and strict orders (see also Def. 5 and Def. 10 below).

DEFINITION 4. Let \mathbf{S} be an arbitrary non-empty subset of $\mathcal{M} \times \mathcal{M}$. The relation \mathbf{S}^t defined by $\alpha \mathbf{S}^t \beta$ for $\alpha, \beta \in \mathcal{M}$, if, for some $n \in \mathbb{N} \setminus \{0, 1\}$, $(\gamma_i, \gamma_{i+1}) \in \mathbf{S}$ for $i = 1, \dots, n-1$ and $\alpha = \gamma_1$ and $\beta = \gamma_n$, is called the transitive closure of \mathbf{S} .

I omit the proof for the following textbook result (cf. [8], [11]).

PROPOSITION 2. The transitive relation generated by \mathbf{S} is its transitive closure,

$$\preceq_{\mathbf{S}} = \mathbf{S}^t. \quad (1)$$

DEFINITION 5. A preorder is a reflexive transitive relation (1 and 4 of Def. 1).

Historically, for instance in the important publication [5] by McNeille, preorders were also simply called ‘orders’. In the contemporary literature (cf. [8]), partial orders are defined as antisymmetric preorders (1, 4, and 6 of Def. 1). The use of the term quasiorder for a preorder seems to be less common (cf. [8], [11]).

DEFINITION 6. A vector preorder is an additive, positively homogeneous preorder (1, 4, 7, and 8 of Def. 1).

LEMMA 1. A total transitive relation (1 and 9 of Def. 1) is reflexive and, thus, a total preorder.

DEFINITION 7 (Reflexive closure). Let \mathbf{S} be an arbitrary non-empty subset of $\mathcal{M} \times \mathcal{M}$. The relation $\mathbf{S}^r = \mathbf{S} \cup \{(\alpha, \alpha) : \alpha \in \mathcal{M}\}$ is called the reflexive closure of \mathbf{S} .

By definition, \mathbf{S}^r is the smallest reflexive relation containing \mathbf{S} . It is straightforward to see that the generation of the transitive closure and of the reflexive closure is commutative, i.e. $(\mathbf{S}^r)^t = (\mathbf{S}^t)^r$. Thus, $\preceq_{\mathbf{S}}^r$ is the uniquely determined smallest preorder which contains \mathbf{S} .

LEMMA 2. Let \preceq_1 and \preceq_2 be two vector preorder. Then, $\preceq_{\preceq_1 \cup \preceq_2}^{add,ph}$ is a vector preorder, and $\preceq_{\preceq_1 \cup \preceq_2}^{add,ph} = \preceq_{\preceq_1 \cup \preceq_2}$.

Therefore, with Prop. 2, the transitive additive positively homogeneous relation generated by two vector preorders is the transitive closure of their union.

Proof of Lemma 2. The relation $\preceq_{\preceq_1 \cup \preceq_2}$ is a preorder, because it is transitive and, with \preceq_1 and \preceq_2 , it again is reflexive. It now suffices to show that $\preceq_{\preceq_1 \cup \preceq_2}$ is additive and positively homogeneous. Positive homogeneity follows via Prop. 2 from Def. 4, because the chain in that definition can be scaled by any $r \in \mathbb{R}_0^+$. Assume now $(\alpha, \beta), (\bar{\alpha}, \bar{\beta}) \in \preceq_{\preceq_1 \cup \preceq_2}$. By Def. 4 and Prop. 2, $(\alpha, \beta) \in \preceq_{\preceq_1 \cup \preceq_2}$ means that for some $n \in \mathbb{N} \setminus \{0, 1\}$, $(\gamma_i, \gamma_{i+1}) \in (\preceq_1 \cup \preceq_2)$ for $i = 1, \dots, n-1$ and $\alpha = \gamma_1$ and $\beta = \gamma_n$. I can add $(\bar{\alpha}, \bar{\alpha})$ to each link of this chain, because of reflexivity and additivity of both, \preceq_1 and \preceq_2 . The so obtained chain with links in $\preceq_1 \cup \preceq_2$ therefore starts with $(\alpha + \bar{\alpha}, \gamma_2 + \bar{\alpha})$ and ends with $(\gamma_{n-1} + \bar{\alpha}, \beta + \bar{\alpha})$. Similarly, there exists a chain for $(\bar{\alpha}, \bar{\beta}) \in \preceq_{\preceq_1 \cup \preceq_2}$ that starts with $(\beta + \bar{\alpha}, \beta + \bar{\gamma}_2)$ and ends with $(\beta + \bar{\gamma}_{\bar{n}-1}, \beta + \bar{\beta})$ for appropriate \bar{n} and $(\bar{\gamma}_i, \bar{\gamma}_{i+1}) \in (\preceq_1 \cup \preceq_2)$ for $i = 1, \dots, \bar{n}-1$. Thus, summarily, a chain with links in $\preceq_1 \cup \preceq_2$ exists that starts with $(\alpha + \bar{\alpha}, \gamma_2 + \bar{\alpha})$ and ends with $(\beta + \bar{\gamma}_{\bar{n}-1}, \beta + \bar{\beta})$. Therefore, $(\alpha + \bar{\alpha}, \beta + \bar{\beta}) \in \preceq_{\preceq_1 \cup \preceq_2}$, and additivity is established. \square

3 Equivalence relations and strict partial orders

In this section, and next to providing more definitions, I derive several results on calculus with relations, as well as some statements concerning the splitting and the merging of transitive relations.

DEFINITION 8. For any relation \mathbf{R} ,

$$\vec{\mathbf{R}} = \{(\alpha, \beta) \in \mathbf{R} : \alpha \mathbf{R} \beta \text{ and } \beta \mathbf{R} \alpha\} \quad (2)$$

is its symmetric part, and

$$\bar{\mathbf{R}} = \mathbf{R} \setminus \vec{\mathbf{R}} \quad (3)$$

its asymmetric part. For a transitive relation \preceq , I denote $\sim = \succcurlyeq$ and $\prec = \preccurlyeq$.

Obviously, one of the parts (2) and (3) can be empty.

Formulated for total transitive relations, Prop. 1.B.1 in [6] implies the first statement of the following lemma.

LEMMA 3. Let \preceq be a transitive relation.

1. If $\alpha \sim \beta$ and $\beta \sim \gamma$, then $\alpha \sim \gamma$. If (a) $\alpha \prec \beta$ and $\beta \sim \gamma$, or if (b) $\alpha \sim \beta$ and $\beta \prec \gamma$, or if (c) $\alpha \prec \beta$ and $\beta \prec \gamma$, then $\alpha \prec \gamma$.
2. Let \preceq also be reflexive and additive. If $\alpha \sim \beta$ and $\gamma \sim \delta$, then $\alpha + \gamma \sim \beta + \delta$. If $\alpha \prec \beta$ and $\gamma \sim \delta$, or if $\alpha \prec \beta$ and $\gamma \prec \delta$, then $\alpha + \gamma \prec \beta + \delta$.
3. Let \preceq also be positively homogeneous, and assume $r \in \mathbb{R}_0^+$. If $\alpha \sim \beta$, then $r\alpha \sim r\beta$. If $\alpha \prec \beta$, then $r\alpha \prec r\beta$.

Proof. 1. Transitivity implies for all these cases $\alpha \preceq \gamma$. The first statement follows, because symmetry of \sim implies that the reverse holds, too. For (a), assume $\gamma \preceq \alpha$. Now, $\alpha \preceq \beta \preceq \gamma \preceq \alpha$, which implies the contradiction $\alpha \sim \beta$. Statements (b) and (c) follow in a similar manner.

2. Additivity implies for all three cases $\alpha + \gamma \preceq \beta + \delta$. The first statement holds, because symmetry of \sim implies that $\beta + \delta \preceq \alpha + \gamma$, as well. The second statement follows, because if $\beta + \delta \preceq \alpha + \gamma$ held true, then $\gamma \preceq \delta$, $-\gamma \preceq -\gamma$, and $-\delta \preceq -\delta$ could be added to obtain the contradiction $\beta \preceq \alpha$. The third statement follows in a similar manner.

3. Since $\alpha \preceq \beta$ and $\beta \preceq \alpha$, the first statement follows immediately from positive homogeneity. For the second one, if $r\beta \preceq r\alpha$, then multiplication with $1/r$ would imply the contradiction $\beta \preceq \alpha$. \square

DEFINITION 9. An equivalence relation is a symmetric preorder (1, 2, and 4 of Def. 1).

For the next definition note that asymmetry implies irreflexivity.

DEFINITION 10. A strict partial order is an asymmetric transitive relation (1, 3, and 5 of Def. 1).

A transitive relation on \mathcal{M} splits into an equivalence relation on a subset of \mathcal{M} and a strict partial order on \mathcal{M} . For total transitive relations, this is (i) and (ii) of Prop. 1.B.1 in [6].

LEMMA 4. *For any transitive relation \preceq on a set \mathcal{M} , if non-empty, \sim is a symmetric transitive relation on \mathcal{M} and therefore an equivalence relation on*

$$\mathcal{M}_{\sim} := \{\alpha \in \mathcal{M} : \alpha \sim \beta \text{ for some } \beta \in \mathcal{M}\}. \quad (4)$$

If non-empty, \prec is a strict partial order on \mathcal{M} . If \preceq has the property 6, or simultaneously the properties 4 and 7, or the property 8 of Def. 1, then so do \sim and \prec .

By (3), \sim and \prec are disjoint (mutually exclusive), i.e. $\sim \cap \prec = \emptyset$.

Proof. \sim is symmetric by definition and inherits transitivity (Lemma 3). It therefore is reflexive, symmetric and transitive on \mathcal{M}_{\sim} . Since \sim is the entire symmetric part of \preceq , the rest, \prec , if non-empty, is asymmetric. Transitivity is given by Lemma 3. Optional antisymmetry: While \prec has the property in the sense that it applies to the empty set, \sim directly inherits it from \preceq . Optional reflexivity and additivity, or optional positive homogeneity of \sim and \prec : Lemma 3. \square

COROLLARY 2. *A preorder on \mathcal{M} splits into an equivalence relation on \mathcal{M} and a strict partial order on \mathcal{M} .*

DEFINITION 11. *If non-empty, \sim is called the implied equivalence relation of the transitive relation \preceq , and \prec , if non-empty, is called its implied strict partial order.*

The following notation can be handy:

$$\succ := \{(\beta, \alpha) : (\alpha, \beta) \in \prec\}. \quad (5)$$

LEMMA 5. *An equivalence relation on a subset of \mathcal{M} is a preorder on \mathcal{M} with an empty implied strict partial order.*

LEMMA 6. *For any transitive relation \preceq , the implied strict partial order \prec is a strict partial order on the equivalence classes of the implied equivalence relation, i.e. on \preceq / \sim , in the sense that for any $\alpha \sim \beta$ and $\gamma \sim \delta$, $\alpha \prec \gamma$ implies $\beta \prec \delta$.*

Essentially, this is Theorem 2.10 in [5]. Note that MacNeille in his Definition 2.4 defines $=$ in the same manner as I defined \sim , but in the case of a preorder. From a contemporary perspective, he thus defined equivalence, not equality (which has no definition).

Proof. Lemma 3. \square

PROPOSITION 3. *Let \preceq_1 and \preceq_2 be two transitive relations. If $\prec_2 \subset \prec_1$ and $\sim_1 \subset \sim_2$, then $\preceq := \sim_1 \cup \prec_2$ is a transitive relation with $\sim = \sim_1$ and $\prec = \prec_2$. If \preceq_1 and \preceq_2 both have the property 4, or 6, or simultaneously the properties 4 and 7, or the property 8 of Def. 1, then so does \preceq .*

Proof. For transitivity, distinguish four cases. Case 1: $\alpha \sim_1 \beta$ and $\beta \sim_1 \gamma$. This implies $\alpha \sim_1 \gamma$. Case 2: $\alpha \sim_1 \beta$ and $\beta \prec_2 \gamma$. Since also $\alpha \sim_2 \beta$, Lemma 3 implies $\alpha \prec_2 \gamma$. Case 3: $\alpha \prec_2 \beta$ and $\beta \sim_1 \gamma$. Since also $\beta \sim_2 \gamma$, Lemma 3 implies $\alpha \prec_2 \gamma$. Case 4: $\alpha \prec_2 \beta$ and $\beta \prec_2 \gamma$. Lemma 3 implies $\alpha \prec_2 \gamma$. Thus, $\alpha \preceq \beta$ and $\beta \preceq \gamma$ implies $\alpha \preceq \gamma$. Further, $\sim = \sim_1$ and $\prec = \prec_2$ must hold since $\sim_1 \cap \prec_2 = \emptyset$. Optional reflexivity follows for \preceq since \sim_1 is reflexive if \preceq_1 and \preceq_2 are. Optional antisymmetry of \preceq is inherited from \preceq_1 . If \preceq_1 and \preceq_2 are additive and reflexive, then a case distinction shows that Lemma 3 and $\sim_1 \subset \sim_2$ imply these properties for \preceq . Lemma 3 also takes care of positive homogeneity. \square

PROPOSITION 4. *Let $\preceq_1 \subsetneq \preceq_2$ be two total preorders. Then $\prec_2 \subsetneq \prec_1$ and $\sim_1 \subsetneq \sim_2$.*

Proof. If $\alpha \preceq_2 \beta$ and $\beta \not\preceq_2 \alpha$, then $\alpha \preceq_1 \beta$ and $\beta \not\preceq_1 \alpha$ follows from the inclusion. Therefore, $\prec_2 \subset \prec_1$. If $\alpha \preceq_1 \beta$ and $\beta \preceq_1 \alpha$, then $\alpha \preceq_2 \beta$ and $\beta \preceq_2 \alpha$ again follows from the inclusion. Therefore, $\sim_1 \subset \sim_2$. Totality implies for $i = 1, 2$

$$\sim_i = \mathcal{M} \times \mathcal{M} \setminus (\prec_i \cup \succ_i). \quad (6)$$

Hence, $\prec_1 = \prec_2$ would imply $\sim_1 = \sim_2$, and thus $\preceq_1 = \preceq_2$, which would be a contradiction. Therefore, with the first part, $\prec_2 \subsetneq \prec_1$, and with (6), $\sim_1 \subsetneq \sim_2$. \square

COROLLARY 3. *Let $\preceq_1 \subsetneq \preceq_2$ be two total preorders. Then $\preceq := \sim_1 \cup \prec_2$ is a preorder. If \preceq_1 and \preceq_2 both have the property 6, 7, or 8 of Def. 1, then so does \preceq .*

Proof. Prop. 4 and Prop. 3. \square

DEFINITION 12. *A partial order is an antisymmetric preorder (1, 4, and 6 of Def. 1).*

LEMMA 7. *Let \preceq be a partial order. Then $\alpha \sim \beta$ if and only if $\alpha = \beta$.*

LEMMA 8. *The reflexive closure $(\prec')^r$ of a strict partial order \prec' is a partial order and the smallest preorder that contains \prec' . Hence, any strict partial order can be constructed by removing the reflexive pairs $\{(\alpha, \alpha) : \alpha \in \mathcal{M}\}$ from a partial order.*

4 Chain consistency

Below, I introduce the for this work fundamental notion of consistency, or chain consistency, between two transitive relations. Examples of consistency and inconsistency are given in this section, and I illustrate why the simple assumption of disjoint symmetric and asymmetric parts across two transitive relations generally is not sufficient for their consistency.

DEFINITION 13. *Two transitive relations \preceq_1 and \preceq_2 are chain-consistent with one another if there exist no $(\gamma_i, \gamma_{i+1}) \in (\preceq_1 \cup \preceq_2)$ for $i = 1, \dots, n-1$ and for some $n \in \mathbb{N} \setminus \{0, 1\}$ such that $\gamma_1 = \gamma_n$ and at least one $(\gamma_i, \gamma_{i+1}) \in \prec_1 \cup \prec_2$.*

Chain consistency means that there is no closed chain made up of \sim_1 and \sim_2 , and at least one of \prec_1 or \prec_2 . For instance, some $\alpha \preceq_1 \beta \prec_2 \gamma \preceq_1 \delta \preceq_1 \alpha$ or some $\alpha \sim_1 \beta \prec_2 \gamma \sim_1 \delta \sim_2 \epsilon \prec_1 \alpha$ would constitute a circular chain as in Def. 13. Chain consistency e.g. prevents that $\alpha \preceq_1 \beta$ while $\beta \prec_2 \alpha$.

EXAMPLE 1. Consider $\mathbb{R} \times \mathbb{R}$ with $(a_1, a_2) \sim' (b_1, b_2)$ if and only if $a_1 - 2a_2 = b_1 - 2b_2$. Moreover, \prec'' being defined as Pareto-better, i.e. $(a_1, a_2) \prec'' (b_1, b_2)$ if and only if either $a_1 \leq a_2$ or $b_1 < b_2$, or $a_1 < a_2$ and $b_1 \leq b_2$, or both. It is easy to check that \prec'' is a strict partial order and that \sim' is an equivalence order. Since $(1, 2) \prec'' (5, 4) \sim' (1, 2)$, the relations \sim' and \prec'' are not chain-consistent with one another.

LEMMA 9. For a transitive relation \preceq , the following are chain-consistent with one another: (1) \preceq and \sim , (2) \preceq and \prec , as well as (3) \sim and \prec . Thus, if two transitive relations are chain-inconsistent, closed chains of the form specified in Def. 13 will have members from both relations.

Proof. For these pairs of relations, any closed chain as in Def. 13 would by Lemma 3 imply $\gamma_1 \prec \gamma_1$, thus contradicting the irreflexivity of \prec . \square

EXAMPLE 2. For a linear space \mathcal{M} , consider a real linear map $f : \mathcal{V} \rightarrow \mathbb{R}$ on a subspace $\mathcal{V} \subset \mathcal{M}$ which is not identical zero. For $\alpha, \beta \in \mathcal{M}$, define $\alpha \preceq_f \beta$ whenever $f(\alpha) \leq f(\beta)$. It is now easy to check that \preceq_f is a vector preorder on \mathcal{M} , and that $\alpha \sim_f \beta$ if and only if $\alpha - \beta \in \text{Ker}(f)$, and, thus, $\alpha \prec_f \beta$ if and only if $f(\alpha) < f(\beta)$. By Lemma 9, \sim_f and \prec_f are consistent with one another.

PROPOSITION 5. If \preceq_1 and \preceq_2 are chain-consistent with one another, then

$$(\sim_1 \cup \sim_2) \cap (\prec_1 \cup \prec_2) = \emptyset. \quad (7)$$

However, in general, (7) does not imply chain consistency of \preceq_1 and \preceq_2 .

Proof. If (7) is violated, then there either exists $\alpha \sim_1 \beta$ with $\alpha \prec_2 \beta$, or $\alpha \sim_2 \beta$ with $\alpha \prec_1 \beta$, both leading to chains excluded by Def. 13. For the second part, consider distinct $\gamma_i, i = 1, \dots, 4$, and

$$\preceq_1 = \sim_1 := \{(\gamma_i, \gamma_j) : (i, j) \in \{(1, 4), (4, 1), (2, 3), (3, 2), (1, 1), (2, 2), (3, 3), (4, 4)\}\}, \quad (8)$$

$$\preceq_2 = \prec_2 := \{(\gamma_1, \gamma_3), (\gamma_2, \gamma_4)\}. \quad (9)$$

It is easy to check that the first is an equivalence relation and the second a strict partial order, and that (7) holds. However, $\gamma_1 \prec_2 \gamma_3 \sim_1 \gamma_2 \prec_2 \gamma_4 \sim_1 \gamma_1$. \square

So, while for a transitive relation \preceq it holds that $\sim \cap \prec = \emptyset$, the reverse is not generally true, i.e. if $\sim_1 \cap \prec_2 = \emptyset$ for an equivalence relation \sim_1 and a strict partial order \prec_2 , then $\sim_1 \cup \prec_2$ is not necessarily a transitive relation.

I now arrive at a first simple result. Two transitive relations are chain-consistent with one another if and only if their transitive closure $\preceq_{\preceq_1 \cup \preceq_2}$ does not contradict either of them.

LEMMA 10. Two transitive relations \preceq_1 and \preceq_2 are chain-consistent with one another if and only if there exist no $\alpha, \beta \in \mathcal{M}$ with $\alpha \preceq_{\preceq_1 \cup \preceq_2} \beta$ such that either $\beta \prec_1 \alpha$, or $\beta \prec_2 \alpha$, or both.

Proof. ‘ \Rightarrow ’: By Def. 4 (recall Prop. 2), if the non-permitted situation of Lemma 10 prevails, then a chain of \preceq_1 and \preceq_2 exists from α to β , and this chain can be closed with $\beta \prec_1 \alpha$ or $\beta \prec_2 \alpha$, such that there is no chain-consistency according to Def. 13.

‘ \Leftarrow ’: If the relations are inconsistent with one another, then there exists a closed chain of \preceq_1 and \preceq_2 as not permitted by Def. 13, where at least one $(\gamma_i, \gamma_{i+1}) \in \prec_1 \cup \prec_2$. Setting $\beta = \gamma_i$ and $\alpha = \gamma_{i+1}$ and using Def. 4, it is now clear that the non-permitted situation of Lemma 10 is established. \square

Obviously, $\sim_1 \cup \sim_2 \subset \sim_{\preceq_1 \cup \preceq_2}$ holds in the situation of Lemma 10. Moreover, the second condition in Lemma 10 can easily be shown to be equivalent to $\prec_1 \cup \prec_2 \subset \prec_{\preceq_1 \cup \preceq_2}$. This brings me to the definition of consistent extensions of transitive relations in the next section.

5 Consistent extensions and completions

This section provides the critical transfinite induction results needed in the proof of my main theorem in the next section. The notion of a consistent extension and of a consistent completion of a transitive relation is introduced, and a technique to extend – and, consequently, complete – a non-total (non-complete) transitive relation in a consistent manner is explained.

The following extension of a relation is structure preserving regarding the symmetric and asymmetric parts of the relation.

DEFINITION 14. *Relation \mathbf{R}_2 is a consistent extension of relation \mathbf{R}_1 , written $\mathbf{R}_1 \sqsubset \mathbf{R}_2$, if $\bar{\mathbf{R}}_1 \subset \bar{\mathbf{R}}_2$ and $\vec{\mathbf{R}}_1 \subset \vec{\mathbf{R}}_2$.*

LEMMA 11. *Let \mathbf{S} be an arbitrary non-empty subset of $\mathcal{M} \times \mathcal{M}$. Then, $\preceq_{\mathbf{S}} \sqsubset \preceq_{\mathbf{S}}^r$.*

LEMMA 12. *If for two transitive relations $\preceq_1 \sqsubset \preceq_2$, then \preceq_1 and \preceq_2 are chain-consistent.*

Proof. Because of the inclusions in Def. 14, chain inconsistency of \preceq_1 and \preceq_2 would imply that of \sim_2 and \prec_2 , which is impossible by Lemma 9. \square

LEMMA 13. *If $\preceq_1 \sqsubset \preceq_2$ for a preorder \preceq_1 and a transitive relation \preceq_2 , then \preceq_2 is a preorder.*

PROPOSITION 6. \sqsubset *is a partial order on all relations on \mathcal{M} which share the same subset of properties from Def. 1. Any totally \sqsubset -ordered subset \mathcal{R} of these relations on \mathcal{M} has the upper bound $\mathbf{B} = \bigcup_{\mathbf{R} \in \mathcal{R}} \mathbf{R}$, which also has the properties. Moreover, $\vec{\mathbf{B}} = \bigcup_{\mathbf{R} \in \mathcal{R}} \vec{\mathbf{R}}$ and $\bar{\mathbf{B}} = \bigcup_{\mathbf{R} \in \mathcal{R}} \bar{\mathbf{R}}$.*

Proof. Clearly, \sqsubset is reflexive, antisymmetric, and transitive, and therefore a partial order on all relations on \mathcal{M} that share the same set of properties from Def. 1. Clearly, $\mathbf{B} = (\bigcup_{\mathbf{R} \in \mathcal{R}} \bar{\mathbf{R}}) \cup (\bigcup_{\mathbf{R} \in \mathcal{R}} \vec{\mathbf{R}})$, and $\bigcup_{\mathbf{R} \in \mathcal{R}} \vec{\mathbf{R}}$ is symmetric. $\bigcup_{\mathbf{R} \in \mathcal{R}} \bar{\mathbf{R}}$ is asymmetric, because, by construction, it cannot have elements of the form (α, α) , and if for $\alpha \neq \beta$ it held that $(\alpha, \beta), (\beta, \alpha) \in \bigcup_{\mathbf{R} \in \mathcal{R}} \bar{\mathbf{R}}$, then, given the total ordering, there would have to be one

member $\vec{\mathbf{R}}$ containing both of them, which cannot be. Thus, $\vec{\mathbf{B}}$ and $\vec{\mathbf{B}}$ are as stated, and $\mathbf{R} \sqsubset \mathbf{B}$ for any $\mathbf{R} \in \mathcal{R}$. Regarding the properties, consider first transitivity. $\alpha \mathbf{R}_1 \beta$ and $\beta \mathbf{R}_2 \gamma$ means that there exist $\mathbf{R}_1, \mathbf{R}_2 \in \mathcal{R}$ such that $\alpha \mathbf{R}_1 \beta$ and $\beta \mathbf{R}_2 \gamma$. If I assume without loss of generality that $\mathbf{R}_1 \sqsubset \mathbf{R}_2$, then also $\alpha \mathbf{R}_2 \beta$, and therefore $\alpha \mathbf{R}_2 \gamma$ by transitivity of \mathbf{R}_2 , which in turn implies $\alpha \mathbf{R} \gamma$. Thus, \mathbf{R} is a transitive relation. With similar or even simpler arguments, the remaining properties 2 to 9 of Def. 1 can be checked to hold for \mathbf{B} if they hold for each member of \mathcal{R} . For instance, for property 3 (asymmetry), if all members of \mathcal{R} have the property, then $\mathbf{B} = \vec{\mathbf{B}}$, so \mathbf{B} has the property, as well. \square

In the following, $\preceq^{(add,ph)}$ means that \preceq applies in cases where only transitivity is considered, and $\preceq^{add,ph}$ applies in cases where, additionally, additivity and positively homogeneity are considered.

PROPOSITION 7. *For $\alpha \neq \beta$, $\alpha, \beta \in \mathcal{M}$, assume that neither $\alpha \preceq \beta$, nor $\beta \preceq \alpha$, for a given transitive relation \preceq .*

1. *If \preceq is a transitive relation, preorder, or strict partial order on the set \mathcal{M} , then*

$$\preceq \sqsubset \preceq_{\preceq \cup \{\alpha, \beta\}} = (\preceq \cup \{(\gamma, \delta) : \gamma \preceq \alpha \text{ or } \gamma = \alpha, \text{ and } \beta \preceq \delta \text{ or } \beta = \delta\}), \quad (10)$$

and the transitive closure $\preceq_{\preceq \cup \{\alpha, \beta\}}$ is again a transitive relation, preorder, or strict partial order.

2. *If \preceq is a vector preorder on the real linear space \mathcal{M} , then*

$$\preceq \sqsubset \preceq_{\preceq \cup \{\alpha, \beta\}}^{add,ph} = (\preceq \cup \{(r\gamma + \epsilon, r\delta + \zeta) : \gamma \preceq \alpha, \beta \preceq \delta, \epsilon \preceq \zeta, r \in \mathbb{R}_{>0}^+\}), \quad (11)$$

where $\preceq_{\preceq \cup \{\alpha, \beta\}}^{add,ph}$ is a vector preorder.

In both cases, $\preceq \neq \preceq_{\preceq \cup \{\alpha, \beta\}}^{(add,ph)}$, and

$$\alpha \prec_{\preceq \cup \{\alpha, \beta\}}^{(add,ph)} \beta, \quad (12)$$

as well as

$$\sim_{\preceq \cup \{\alpha, \beta\}}^{(add,ph)} = \sim. \quad (13)$$

This means that the transitive relation, preorder, or strict partial order $\preceq_{\preceq \cup \{\alpha, \beta\}}$ is the smallest one larger than \preceq that also strictly prefers β over α . Moreover, only the strict part of \preceq is extended by this procedure. A result in the case of partial orders that is somewhat related to Proposition 7 is Theorem 1.19 in [8].

Proof. Part 1. Denote the right hand side of (10) short with \preceq' . Note that $\alpha \preceq' \beta$. Clearly, $\preceq \subset \preceq'$, and \preceq' is reflexive in the preorder case. For the second equality in (10), it only needs to be shown that the relation on the right, which is \preceq' , is transitive, since it obviously is contained in $\preceq_{\preceq \cup \{\alpha, \beta\}}$. For transitivity, assume $\gamma_1 \preceq' \gamma_2$ and $\gamma_2 \preceq' \gamma_3$. Case 1: $\gamma_1 \preceq \gamma_2$ and $\gamma_2 \preceq \gamma_3$. Then, $\gamma_1 \preceq \gamma_3$, and thus $\gamma_1 \preceq' \gamma_3$. Case 2: $\gamma_1 \preceq \gamma_2$, but $\gamma_2 \not\preceq \gamma_3$. Then, $\gamma_1 \preceq \gamma_2 \preceq \alpha$ or $\gamma_1 \preceq \gamma_2 = \alpha$, and $\beta \preceq \gamma_3$ or $\beta = \gamma_3$, and thus, because $\gamma_1 \preceq \alpha$,

$\gamma_1 \preceq' \gamma_3$. Case 3: $\gamma_1 \not\preceq \gamma_2$, but $\gamma_2 \preceq \gamma_3$. Then, $\gamma_1 \preceq \alpha$ or $\gamma_1 = \alpha$, and $\beta \preceq \gamma_2 \preceq \gamma_3$ or $\beta = \gamma_2 \preceq \gamma_3$, and thus, because $\beta \preceq \gamma_3$, $\gamma_1 \preceq' \gamma_3$. Case 4: $\gamma_1 \not\preceq \gamma_2$ and $\gamma_2 \not\preceq \gamma_3$. Then, $\gamma_1 \preceq \alpha$ or $\gamma_1 = \alpha$, and $\beta \preceq \gamma_2$ or $\beta = \gamma_2$, as well as $\gamma_2 \preceq \alpha$ or $\gamma_2 = \alpha$, and $\beta \preceq \gamma_3$ or $\beta = \gamma_3$. This implies $\beta \preceq \alpha$ or $\beta = \alpha$, which is a contradiction to the assumptions. Thus, this case cannot occur. It is now established that the right hand side of (10) is a transitive relation, and a preorder, if \preceq is one. Since $\beta \preceq' \alpha$ would with (10) either imply $\alpha = \beta$, or imply both, $\alpha \preceq \beta$ and $\beta \preceq \alpha$, which is a contradiction to the initial assumptions, $\beta \not\preceq' \alpha$, and (12) holds. This, in turn, implies via Lemma 3 and the right hand side of (10) that $\sim' = \sim$, and (13) holds. Since $\preceq \subset \preceq'$, $\sim' = \sim$ implies $\prec \subset \prec'$, and, in conclusion, \square holds on the left hand side in (10). However, with Lemma 4, this also implies that \preceq' is a strict partial order if \preceq is, since then $\sim' = \sim = \emptyset$.

Part 2. Again, denote the right hand side of (11) short with \preceq' . Reflexivity of the second and third relation in (11) is given by the reflexivity of \preceq . For the second equality in (11), it now only needs to be shown that the relation on the right is additive, and positively homogeneous, because reflexivity and additivity imply transitivity, and thus a preorder, and it is easy to see that the relation on the right must be contained in $\preceq_{\preceq \cup \{(\alpha, \beta)\}}^{add, ph}$, which, since in this case \preceq is reflexive, is the transitive, reflexive, additive, and positively homogeneous relation generated by $\preceq \cup \{(\alpha, \beta)\}$. For additivity, first assume $r_1\gamma_1 + \epsilon_1 \preceq' r_1\delta_1 + \zeta_1$ and $r_2\gamma_2 + \epsilon_2 \preceq' r_2\delta_2 + \zeta_2$ as in the right hand side of (11). Additivity and positive homogeneity of \preceq imply $(\epsilon_1 + \epsilon_2) \preceq (\zeta_1 + \zeta_2)$, $(r_1\gamma_1 + r_2\gamma_2)/(r_1 + r_2) \preceq \alpha$, and $\beta \preceq (r_1\delta_1 + r_2\delta_2)/(r_1 + r_2)$. Therefore, (11) implies

$$\epsilon_1 + \epsilon_2 + r_1\gamma_1 + r_2\gamma_2 \preceq' \zeta_1 + \zeta_2 + r_1\delta_1 + r_2\delta_2. \quad (14)$$

Assume now $r_1\gamma + \epsilon_1 \preceq' r_1\delta + \zeta_1$ as in the right hand side of (11), as well as $\epsilon_2 \preceq \zeta_2$. Since $\epsilon_1 \preceq \zeta_1$, and therefore $\epsilon_1 + \epsilon_2 \preceq \zeta_1 + \zeta_2$, and since $\gamma \preceq \alpha$ and $\beta \preceq \delta$, (11) implies

$$\epsilon_1 + \epsilon_2 + r_1\gamma \preceq' \zeta_1 + \zeta_2 + r_1\delta. \quad (15)$$

Therefore, additivity is established. The positive homogeneity of the right hand side of (11) is straightforward to see, thus, \preceq' is an additive positively homogeneous preorder. Assume now $\beta \preceq' \alpha$. According to (11), and since $\beta \not\preceq \alpha$, there exist $\gamma \preceq \alpha$, $\beta \preceq \delta$, $\epsilon \preceq \zeta$, and $r \in \mathbb{R}_{>0}^+$ such that $\beta = r\gamma + \epsilon$ and $\alpha = r\delta + \zeta$. This implies $\gamma = (\beta - \epsilon)/r$ and $\delta = (\alpha - \zeta)/r$. Since $\gamma \preceq \alpha$ and $\beta \preceq \delta$, one obtains $(\beta - \epsilon)/r \preceq \alpha$ and $\beta \preceq (\alpha - \zeta)/r$. Multiplying each with r and adding them,

$$(1+r)\beta - \epsilon \preceq (1+r)\alpha - \zeta. \quad (16)$$

Adding $\epsilon \preceq \zeta$ and multiplying with $1/(1+r)$, I obtain the contradiction $\beta \preceq \alpha$. Thus, (12) holds. This, in turn, implies via Lemma 3 and the right hand side of (11) that $\sim' = \sim$, and (13) holds. \square

COROLLARY 4. *If a transitive relation \preceq_1 and a strict partial order \prec_2 are chain-consistent with one another, then the implied equivalence relation \sim_1 and \prec_2 are chain-consistent with one another. The reverse is not necessarily true.*

Proof. If \sim_1 and \prec_2 are not chain-consistent, there is a circular chain as described in Def. 13 for \sim_1 and \prec_2 . However, since $\sim_1 \subset \preceq_1$, this implies a circular chain as in Def. 13

for \preceq_1 and \prec_2 . To see that the reverse is not true, use for a non-total equivalence relation, \sim , Prop. 7, (10), with some $\alpha \not\sim \beta$ to obtain

$$\preceq_1 := \preceq_{\sim \cup \{(\alpha, \beta)\}} = \sim \cup \{(\gamma, \delta) : \gamma \sim \alpha \text{ or } \gamma = \alpha, \text{ and } \beta \sim \delta \text{ or } \beta = \delta\}, \quad (17)$$

$$\preceq_2 := \preceq_{\sim \cup \{(\beta, \alpha)\}} = \sim \cup \{(\gamma, \delta) : \gamma \sim \beta \text{ or } \gamma = \beta, \text{ and } \alpha \sim \delta \text{ or } \alpha = \delta\}. \quad (18)$$

By Prop. 7, Eq. (13), $\sim_1 = \sim_2 = \sim$, which implies with Lemma 9 that \sim_1 is consistent with \prec_2 , but $\alpha \preceq_1 \beta \prec_2 \alpha$, so \preceq_1 and \prec_2 are not chain-consistent with one another. \square

DEFINITION 15. *For a relation \mathbf{R} on \mathcal{M} , the relation $\overline{\mathbf{R}}$ on \mathcal{M} is a consistent completion of \mathbf{R} if $\overline{\mathbf{R}}$ is total and $\mathbf{R} \sqsubset \overline{\mathbf{R}}$.*

Note that any relation is trivially extended to a total one by $\mathcal{M} \times \mathcal{M}$. However, structure-preserving, i.e. consistent completion is less trivial. Recall from Lemma 1 that a total transitive relation is a preorder.

THEOREM 1. *Any transitive relation has a consistent completion which is a preorder. Any vector preorder has a consistent completion which again is a vector preorder.*

Proof. Let \preceq be a non-total transitive relation or vector preorder on a set or real linear space \mathcal{M} . Consider now the set of all transitive relations or all vectors preorders on \mathcal{M} which are at least as large as \preceq with respect to \sqsubset . Clearly, \sqsubset is a partial order on these relations. By Prop. 6, any totally ordered subset has an upper bound. Therefore, by Zorn's Lemma (cf. [12]; the usage of Zorn's "maximum principle" here is exactly as formulated by him, namely for inclusion orders on sets), these relations have at least one maximal element, of which one is chosen and denoted by $\overline{\preceq}$. Assume now that for some $\alpha, \beta \in \mathcal{M}$ neither $\alpha \overline{\preceq} \beta$, nor $\beta \overline{\preceq} \alpha$. If $\alpha \neq \beta$, then, by Prop. 7, $\preceq_{\overline{\preceq} \cup \{(\alpha, \beta)\}}^{(add, ph)}$ is strictly \sqsubset -larger than $\overline{\preceq}$, which is a contradiction to $\overline{\preceq}$ being a maximal element. In the case of transitive non-reflexive relations, $\alpha = \beta$ is a possibility, too. Then it is straightforward to see that $\overline{\preceq} \cup \{(\alpha, \alpha)\}$ is strictly \sqsubset -larger than $\overline{\preceq}$ and transitive, which, again, is a contradiction to $\overline{\preceq}$ being a maximal element. \square

6 A consistency theorem for transitive relations

This section proves my main theorem and states two of its corollaries.

THEOREM 2. *Let \preceq_1 and \preceq_2 be transitive relations on a set \mathcal{M} .*

1. *The following are equivalent:*

- (a) \preceq_1 and \preceq_2 are chain-consistent with one another.
- (b) $\preceq_{\preceq_1 \cup \preceq_2}$ is chain-consistent with \preceq_1 and with \preceq_2 .
- (c) $\preceq_{\preceq_1 \cup \preceq_2}$ consistently extends \preceq_1 and \preceq_2 .
- (d) There exists a transitive relation that consistently extends \preceq_1 and \preceq_2 .
- (e) There exists a common transitive consistent completion for \preceq_1 and \preceq_2 .

(f) There exists a transitive consistent completion $\overline{\preceq}_1$ of \preceq_1 , and $\overline{\preceq}_1$ and \preceq_2 are chain-consistent with one another.

2. \preceq_1 and \preceq_2 are chain-consistent with one another and $\preceq_{\preceq_1 \cup \preceq_2}^r$ is total if and only if there exists a uniquely determined common transitive consistent completion for \preceq_1 and \preceq_2 . In this case, the preorder $\preceq_{\preceq_1 \cup \preceq_2}^r$ is the consistent completion.
3. If \preceq_1 and \preceq_2 are vector preorders on the real linear space \mathcal{M} , the statements above hold while $\preceq_{\preceq_1 \cup \preceq_2}$ and the extensions, completions, or total preorders again are vector preorders.

Proof. Part 1. ‘(a) \Rightarrow (b)’: Assume there is a closed chain for $\preceq_{\preceq_1 \cup \preceq_2}$ and \preceq_2 as in Def. 13. Since, by Prop. 2, $\preceq_{\preceq_1 \cup \preceq_2}$ is the transitive closure of \preceq_1 and \preceq_2 , there is a contradiction to (a) if the chain contains at least one \prec_2 . Furthermore, it is straightforward to see that $\sim_2 \subset \sim_{\preceq_1 \cup \preceq_2}$. Thus, the chain must lie entirely in $\preceq_{\preceq_1 \cup \preceq_2}$ with at least one link from $\prec_{\preceq_1 \cup \preceq_2}$, but this contradicts Lemma 9 (3). A similar argument proves chain consistency with \preceq_1 .

‘(b) \Rightarrow (c)’: It is easy to see that $(\sim_1 \cup \sim_2) \subset \sim_{\preceq_1 \cup \preceq_2}$. Assuming that (c) does not hold, $(\prec_1 \cup \prec_2) \not\subset \prec_{\preceq_1 \cup \preceq_2}$ follows. Then at least one of the two following cases applies. Case 1: $\alpha \prec_2 \beta$ and $\alpha \sim_{\preceq_1 \cup \preceq_2} \beta$, where the latter implies $\beta \preceq_{\preceq_1 \cup \preceq_2} \alpha$. Therefore, $\alpha \prec_2 \beta \preceq_{\preceq_1 \cup \preceq_2} \alpha$, a contradiction to (b). Case 2: $\alpha \prec_1 \beta$ and $\alpha \sim_{\preceq_1 \cup \preceq_2} \beta$, where the latter implies $\beta \preceq_{\preceq_1 \cup \preceq_2} \alpha$. Therefore, $\alpha \prec_1 \beta \preceq_{\preceq_1 \cup \preceq_2} \alpha$, again a contradiction to (b).

‘(c) \Rightarrow (d)’: Trivial. ‘(d) \Rightarrow (e)’: Theorem 1. ‘(e) \Rightarrow (f)’: Lemma 12.

‘(f) \Rightarrow (a)’: Assume there is a closed chain for \preceq_1 and \preceq_2 as in Def. 13. Because $\preceq_1 \sqsubset \overline{\preceq}_1$, this means that there is a closed chain as in Def. 13 with members from both, $\overline{\preceq}_1$ and \preceq_2 , which is a contradiction to (f).

Part 2. ‘ \Rightarrow ’: With part 1 and with Lemma 11, $\preceq := \preceq_{\preceq_1 \cup \preceq_2}^r$ consistently completes \preceq_1 and \preceq_2 . Assume now that \preceq° is a reflexive transitive relation (i.e. a preorder) different from \preceq , for which – like for \preceq – also $\preceq_1 \sqsubset \preceq^\circ$ and $\preceq_2 \sqsubset \preceq^\circ$. Since \preceq is inclusion-minimal (cf. Def. 2 and Cor. 1) with these properties, $\preceq \subset \preceq^\circ$ holds. Because of this inclusion, \preceq° must be total. Moreover, since \preceq° is supposed to be different from \preceq , one has $\preceq \subsetneq \preceq^\circ$. By Cor. 3 and by Prop. 4, $\preceq_3 := \sim \cup \prec^\circ$ now is a preorder which is strictly inclusion-smaller than \preceq . Furthermore $\preceq_1 \sqsubset \preceq_3$ and $\preceq_2 \sqsubset \preceq_3$. This is a contradiction to the minimality of \preceq .

‘ \Leftarrow ’: Chain consistency of \preceq_1 and \preceq_2 is clear from part 1. It also follows from part 1 and from Lemma 11 that the preorder $\preceq := \preceq_{\preceq_1 \cup \preceq_2}^r$ consistently extends \preceq_1 and \preceq_2 . If \preceq is non-total, and therefore for some $\alpha, \beta \in \mathcal{M}$ with $\alpha \neq \beta$ neither $\alpha \preceq \beta$, nor $\beta \preceq \alpha$, Prop. 7 shows that

$$\preceq \sqsubset \preceq_{\preceq \cup \{(\alpha, \beta)\}}, \quad (19)$$

$$\preceq \sqsubset \preceq_{\preceq \cup \{(\beta, \alpha)\}}, \quad (20)$$

$$\preceq_{\preceq \cup \{(\beta, \alpha)\}} \neq \preceq_{\preceq \cup \{(\alpha, \beta)\}}, \quad (21)$$

$$(\alpha, \beta) \subset \prec_{\preceq \cup \{(\alpha, \beta)\}}, \quad (22)$$

$$(\beta, \alpha) \subset \prec_{\preceq \cup \{(\beta, \alpha)\}}. \quad (23)$$

$\preceq_{\preceq \cup \{(\alpha, \beta)\}}$ and $\preceq_{\preceq \cup \{(\beta, \alpha)\}}$ can be consistently completed to transitive relations by Theorem 1. Since (α, β) will be in the asymmetric part of one of the completions, and (β, α) will be

in the corresponding part of the other one, this proves that the *a priori* assumed common completion is not unique in this case. Therefore, uniqueness of the completion implies totality of \preceq .

Part 3. For part 1 this holds because of Lemma 2 and Theorem 1. For part 2, Lemma 2, Cor. 3, Prop. 4, Prop. 7, and Theorem 1 cover the additive positively homogeneous case. \square

The equivalence of 1 (a) and 1 (c) of Theorem 2 was already implied by Lemma 10, the remarks underneath its proof, and by Def. 14. Beside the vector preorder case and other statements, the theorem now, of course, added the completeness and uniqueness results.

COROLLARY 5. *If \preceq_1 and \preceq_2 are chain-consistent with one another and $\preceq_{\preceq_1 \cup \preceq_2}^r$ is total, then $\preceq_{\preceq_1 \cup \preceq_2}^r$ is the only preorder \preceq fulfilling $\preceq_1 \sqsubset \preceq$ and $\preceq_2 \sqsubset \preceq$.*

Next, I finally provide an explicit statement of the meaningful reverse of the dissection of any preorder into an equivalence relation and a strict partial order that I mentioned in the first section and that was formalized by Lemma 4. If the equivalence relation and the strict partial order are chain-consistent with one another, then their transitive closure consistently extends both of them, meaning: Old equivalence then implies new (extended) equivalence, and old strict ordering implies new (extended) strict ordering. Furthermore, this transitive closure also can be completed while keeping the just explained implications in place.

The following formulation of my result also seems to be related to the Fundamental Theorem of Asset Pricing, which sits at the heart of financial mathematics. The analogies will be explored in Section 8.

COROLLARY 6. *Let \sim' be an equivalence relation and let \prec'' be a strict partial order on a set \mathcal{M} .*

1. *The following are equivalent:*

- (a) \sim' and \prec'' are chain-consistent with one another.
- (b) $\preceq_{\sim' \cup \prec''}$ and \prec'' are chain-consistent with one another.
- (c) For the preorder $\preceq_{\sim' \cup \prec''}$ generated by \sim' and \prec'' ,

$$\sim' \subset \sim_{\sim' \cup \prec''} \quad \text{and} \quad (24)$$

$$\prec'' \subset \prec_{\sim' \cup \prec''}. \quad (25)$$

(d) *There exists a total preorder $\overline{\preceq}$ on \mathcal{M} such that*

$$\sim' \subset \overline{\preceq} \quad \text{and} \quad (26)$$

$$\prec'' \subset \overline{\preceq}. \quad (27)$$

- 2. \sim' and \prec'' are chain-consistent with one another and $\preceq_{\sim' \cup \prec''}$ is total if and only if there exists a uniquely determined total preorder $\overline{\preceq}$ on \mathcal{M} with (26) and (27). In this case, $\overline{\preceq} = \preceq_{\sim' \cup \prec''}$.

3. If \sim' and \prec'' are chain-consistent and $\preceq_{\sim' \cup \prec''}$ is total, then $\preceq_{\sim' \cup \prec''}$ is the only preorder \preceq fulfilling

$$\sim' \subset \sim \quad \text{and} \quad (28)$$

$$\prec'' \subset \prec. \quad (29)$$

Proof. Theorem 2 for $\preceq_1 = \sim'$ and $\preceq_2 = \prec''$. Note that $\preceq_{\sim' \cup \prec''}$ is trivially chain-consistent with \sim' . Also note that $\preceq_{\sim' \cup \prec''}$ is reflexive, because \sim' as an equivalence relation is already. \square

7 Application in geometry: Cones

Let \preceq be a vector preorder on the real linear space \mathcal{M} . Let $\mathcal{C} \subset \mathcal{M}$ be a positive cone with vertex o (the nullvector), i.e. for any $\alpha, \beta \in \mathcal{C}$ and $a, b \in \mathbb{R}_0^+$, $a\alpha + b\beta \in \mathcal{C}$. In the literature, a positive cone as defined here is sometimes called convex, non-negative, or simply a ‘cone’.

DEFINITION 16.

$$\mathcal{C}_{\preceq} = \{\gamma \in \mathcal{M} : o \preceq \gamma\} \quad (30)$$

and

$$\alpha \preceq_{\mathcal{C}} \beta \quad \text{if and only if} \quad \beta - \alpha \in \mathcal{C}. \quad (31)$$

The following statement is well-known (e.g. [7] or [10]).

PROPOSITION 8. \mathcal{C}_{\preceq} is a positive cone in \mathcal{M} and $\preceq_{\mathcal{C}}$ is a vector preorder on \mathcal{M} . Furthermore, $\mathcal{C}_{\preceq} = \mathcal{C}$ if and only if $\preceq = \preceq_{\mathcal{C}}$.

Prop. 8 establishes a bijection between the vector preorders on \mathcal{M} and the positive cones in \mathcal{M} .

It is well-known (e.g. [7]) that $\mathcal{C} \cap (-\mathcal{C})$ is the largest linear space in \mathcal{C} . By (31),

$$o \prec_{\mathcal{C}} \gamma \quad \text{if and only if} \quad \gamma \in \mathcal{C} \setminus (-\mathcal{C}). \quad (32)$$

Therefore, the following definition seems appropriate.

DEFINITION 17.

$$\mathcal{C}^L = \mathcal{C} \cap (-\mathcal{C}) \quad (33)$$

is called the linear part of \mathcal{C} and

$$\mathcal{C}^+ = \mathcal{C} \setminus (-\mathcal{C}) \quad (34)$$

is called the non-linear part or strictly positive part of \mathcal{C} .

LEMMA 14. For any $\alpha \in \mathcal{C}^+$, $\beta \in \mathcal{C}$ and $a, b > 0$, it holds that $a\alpha + b\beta \in \mathcal{C}^+$. Moreover,

$$\alpha \sim_{\mathcal{C}} \beta \quad \text{if and only if} \quad \beta - \alpha \in \mathcal{C}^L, \quad (35)$$

$$\alpha \prec_{\mathcal{C}} \beta \quad \text{if and only if} \quad \beta - \alpha \in \mathcal{C}^+. \quad (36)$$

Proof. $a\alpha + b\beta \in \mathcal{C}$ holds. If also $a\alpha + b\beta \in (-\mathcal{C})$, then, since $-b\beta \in (-\mathcal{C})$ and since $(-\mathcal{C})$ is a positive cone as well, it follows that $a\alpha, \alpha \in (-\mathcal{C})$, which is a contradiction. (35) follows directly from (31) and (33), and (36) is (32). \square

Together with Lemma 7, these observations lead to the following result.

COROLLARY 7. *There exists a bijection between the additive positively homogeneous partial orders (i.e. vector partial orders) on \mathcal{M} and the positive cones in \mathcal{M} for which $\mathcal{C}^L = \{o\}$.*

The 1-to-1 correspondence of Cor. 7 is a restriction of the one from Prop. 8.

DEFINITION 18. *If $\mathcal{C} \cup (-\mathcal{C}) = \mathcal{M}$, I call \mathcal{C} complete or total,*

It is easy to show that this condition holds if and only if $\prec_{\mathcal{C}}$ is total.

COROLLARY 8. *If $\gamma_i \in \mathcal{C}$, $i = 1, \dots, n$, with at least one $k \in \{1, \dots, n\}$ where $\gamma_k \in \mathcal{C}^+$, then $\sum_{i=1}^n \gamma_i \in \mathcal{C}^+$.*

Proof. By induction from Lemma 14. \square

Expressed differently, Corollary 8 states that a walk, where each step is a member of the cone and at least one step is in the strictly positive direction of the cone, cannot return to the same place and, in fact, ends up at a new place in strictly positive direction from the old position.

Using again the analogy of a walk, what if one takes each step from one of two different cones, \mathcal{C}_1 and \mathcal{C}_2 , and makes sure at least one step is from the strictly positive part of one of the cones - is it then possible to return to the same spot? Theorem 2 in its version for vector preorders gives the answer that never being able to return is equivalent to that both cones can be embedded into a larger cone such that both linear parts are embedded into the linear part of the larger cone, and both strictly positive parts are embedded into the strictly positive part. To properly formulate this statement, I need the following definitions.

DEFINITION 19. *Two positive cones \mathcal{C}_1 and \mathcal{C}_2 in a real linear space \mathcal{M} are path-consistent with one another if there exist no $\delta_i \in \mathcal{C}_1 \cup \mathcal{C}_2$, $i = 1, \dots, m$, with at least one $k \in \{1, \dots, m\}$ where $\delta_k \in \mathcal{C}_1^+ \cup \mathcal{C}_2^+$ such that $\sum_{i=1}^m \delta_i = o$.*

DEFINITION 20. *A positive cone \mathcal{C}_2 consistently extends a positive cone \mathcal{C}_1 in a real linear space \mathcal{M} if $\mathcal{C}_1^L \subset \mathcal{C}_2^L$ and $\mathcal{C}_1^+ \subset \mathcal{C}_2^+$. If \mathcal{C}_2 is complete, then it is called a consistent completion of \mathcal{C}_1 .*

DEFINITION 21. *For two subsets A and B of a real linear space \mathcal{M} , I define*

$$A \uplus B = A \cup (A + B) \cup B. \quad (37)$$

Obviously, \uplus only differs from $+$ if $o \notin A$ or $o \notin B$.

COROLLARY 9. *Given two positive cones \mathcal{C}_1 and \mathcal{C}_2 in a real linear space \mathcal{M} , the following are equivalent.*

1. \mathcal{C}_1 and \mathcal{C}_2 are path-consistent with one another.
2. There exist no $\delta_i \in \mathcal{C}_1 + \mathcal{C}_2$, $i = 1, \dots, m$, with at least one $k \in \{1, \dots, m\}$ where $\delta_k \in \mathcal{C}_1^+ \uplus \mathcal{C}_2^+$ such that $\sum_{i=1}^m \delta_i = o$.
3. There exists a cone $\mathcal{C} \subset \mathcal{M}$ that consistently extends \mathcal{C}_1 and \mathcal{C}_2 . \mathcal{C} can be chosen to be a consistent completion.
4. There exists a cone $\mathcal{C} \subset \mathcal{M}$ such that

$$\mathcal{C}_1^L + \mathcal{C}_2^L \subset \mathcal{C}^L \quad \text{resp.} \quad (38)$$

$$\mathcal{C}_1^+ \uplus \mathcal{C}_2^+ \subset \mathcal{C}^+. \quad (39)$$

\mathcal{C} can be chosen to be complete.

5. The vector preorders $\preceq_{\mathcal{C}_1}$ and $\preceq_{\mathcal{C}_2}$ are chain-consistent with one another.

Thus, two positive cones are path-consistent with one another if and only if there exists a common consistent completion of them.

Proof. ‘1 \Rightarrow 2’: If a set of δ_i as in statement 2 did exist (as opposed to not existing, as required), then splitting each δ_i into its constituents from \mathcal{C}_1 and from \mathcal{C}_2 or, if necessary, into its constituents from \mathcal{C}_1^+ and from \mathcal{C}_2^+ would create a set in direct contradiction to statement 1.

‘1 \Leftarrow 2’: This is straightforward since $\mathcal{C}_1 \cup \mathcal{C}_2 \subset \mathcal{C}_1 + \mathcal{C}_2$ and $\mathcal{C}_1^+ \cup \mathcal{C}_2^+ \subset \mathcal{C}_1^+ \uplus \mathcal{C}_2^+$.

‘1 \Rightarrow 5’: Because of Lemma 14, it is easy to check that, for $\preceq_{\mathcal{C}_1}$ and $\preceq_{\mathcal{C}_2}$, a chain with members (γ_i, γ_{i+1}) as in Def. 13 would via $\delta_i = \gamma_{i+1} - \gamma_i$ for $i = 1, \dots, n-1$ and $m = n-1$ imply $\{\delta_1, \dots, \delta_m\}$ as in Def. 19.

‘1 \Leftarrow 5’: Vice versa, because of Lemma 14, $\{\delta_1, \dots, \delta_m\}$ as in Def. 19 would via $\gamma_1 = o$, $\gamma_i = \sum_{k=1}^{i-1} \delta_k$ for $i = 2, \dots, m+1$ and $n = m+1$ imply a chain as in Def. 13 for $\preceq_{\mathcal{C}_1}$ and $\preceq_{\mathcal{C}_2}$, which starts and ends in o .

‘3 \Rightarrow 5’: The third statement implies

$$\mathcal{C}_1^L \cup \mathcal{C}_2^L \subset \mathcal{C}^L \quad \text{and} \quad (40)$$

$$\mathcal{C}_1^+ \cup \mathcal{C}_2^+ \subset \mathcal{C}^+, \quad (41)$$

and therefore implies with Lemma 14 that the vector preorder $\preceq_{\mathcal{C}}$ consistently extends the vector preorders $\preceq_{\mathcal{C}_1}$ and $\preceq_{\mathcal{C}_2}$.

‘3 \Leftarrow 5’: Vice versa, if the vector preorder $\preceq_{\mathcal{C}}$ consistently extends the vector preorders $\preceq_{\mathcal{C}_1}$ and $\preceq_{\mathcal{C}_2}$, then Lemma 14 implies (40) and (41). Regarding completeness in statement 3, the equivalence of the completeness of \mathcal{C} and $\preceq_{\mathcal{C}}$ had already been mentioned. Thus, via the corresponding completeness statement of Theorem 2, equivalence of statement 1 and 3 is given.

‘3 \Rightarrow 4’: Follows directly from the additivity of \mathcal{C}^L and \mathcal{C}^+ .

‘3 \Leftarrow 4’: Follows directly from $\mathcal{C}_1 \cup \mathcal{C}_2 \subset \mathcal{C}_1 + \mathcal{C}_2$ and $\mathcal{C}_1^+ \cup \mathcal{C}_2^+ \subset \mathcal{C}_1^+ \uplus \mathcal{C}_2^+$. \square

Before a closing example for this section, and in analogy to Prop. 5, note that while path consistency of two cones via Corollary 9 implies

$$(\mathcal{C}_1^L \cup \mathcal{C}_2^L) \cap (\mathcal{C}_1^+ \cup \mathcal{C}_2^+) = \emptyset, \quad (42)$$

the reverse generally does not hold. A simple example for this in $\mathcal{M} = \mathbb{R}^2$ is given by $\mathcal{C}_1 = \mathbb{R} \times \mathbb{R}_0^+$ and $\mathcal{C}_2 = \mathbb{R} \times \mathbb{R}_0^-$.

EXAMPLE 3. *Expanding on Example 2 with the vector preorder \preceq_f on \mathcal{M} stemming from a linear functional on a subspace $\mathcal{V} \subset \mathcal{M}$, one obtains from (30) and (33) that*

$$\mathcal{C}_{\preceq_f} = f^{-1}(\mathbb{R}_0^+), \quad (43)$$

$$\mathcal{C}_{\preceq_f}^L = f^{-1}(0) = \text{Ker}(f), \quad \text{and} \quad (44)$$

$$\mathcal{C}_{\preceq_f}^+ = f^{-1}(\mathbb{R}_{>0}^+), \quad (45)$$

where the last one could be empty. For a positive cone $\mathcal{C} \subset \mathcal{M}$, assume now that there exist no $\delta_i \in f^{-1}(\mathbb{R}_0^+) \cup \mathcal{C}$, $i = 1, \dots, m$, with at least one $k \in \{1, \dots, m\}$ where $\delta_k \in f^{-1}(\mathbb{R}_{>0}^+) \cup \mathcal{C}^+$ such that $\sum_{i=1}^m \delta_i = o$. This means that \mathcal{C} and the positive cone \mathcal{C}_{\preceq_f} implied by f are path-consistent. It now holds that

$$\mathcal{V} \cap \mathcal{C} \subset f^{-1}(\mathbb{R}_0^+), \quad (46)$$

$$\mathcal{V} \cap \mathcal{C}^L \subset f^{-1}(0), \quad \text{and} \quad (47)$$

$$\mathcal{V} \cap \mathcal{C}^+ \subset f^{-1}(\mathbb{R}_{>0}^+). \quad (48)$$

To see (46), assume that $f(\alpha) < 0$ for $\alpha \in \mathcal{V} \cap \mathcal{C}$. Then, $\alpha - \alpha = o$, but $\alpha \in \mathcal{C}$ and $-\alpha \in f^{-1}(\mathbb{R}_{>0}^+)$, which is a contradiction to the path consistency requirement. (47) and (48) now follow with (44) and (45) from the consistent common extension statement of Corollary 9.

(47) with (48) is a stricter condition than the sometimes considered mere ‘positivity’, which usually is understood as non-negativity, of f on the cone \mathcal{C} (e.g. [2]). As such, the consistency between \mathcal{C} and f , which is given as the path-consistency between \mathcal{C} and \mathcal{C}_{\preceq_f} , does not only imply a total cone which consistently extends \mathcal{C} and \mathcal{C}_{\preceq_f} , but it also implies a strict type of positivity of f on \mathcal{C} .

In Example 3, if there existed a linear extension F of f onto \mathcal{M} which implied the extending cone, this would be an addition of precision to the M. Riesz Extension Theorem (e.g. [2]). However, I was not able to generally derive such a result.

8 Application in financial economics

In financial economics or mathematical finance, an often encountered model setup is a market of traded financial goods – meaning of deterministic or stochastic cash flows. Typically, exchangeability is indirectly described by the property of having the same price, where price is the multiple of the numéraire, or unit of account, against which a good can be traded, while it is assumed that all *a priori* given goods of the given market can be

traded for a certain amount of the numéraire (for the more mathematical articles on this topic, see [3] and the references therein). Of course, any good can be exchanged with itself. While usually not stated explicitly and because, typically, the price functional is *a priori* assumed to be linear in the goods, the relation of spot trades at time zero established by these assumptions usually creates a linear equivalence relation on the goods, which are typically given in the form of a real linear space. However, below, I will not require any linearity assumptions.

Another part of the model setups is that there is an understanding of what makes one financial good, or cash flow, for all market participants – and, thus, objectively – strictly better than another. For stochastic cash flows, one cash flow usually is considered strictly better than another one if in all economic scenarios, meaning in any combination of a time and of an event with a non-zero probability, it pays at least as much as the corresponding other cash flow, and in at least one such scenario it pays strictly more.

In economic theory, preorders are used as preference orders which describe the economic preferences of market participants regarding certain goods, bundles of goods, or alternatives. Strict partial orders are used as strict preference orders, and equivalence relations emerge as indifference relations. For an introduction to economic preference orders see e.g. [6], a widely known and appreciated textbook. Reference [1] (English translation in: *Preferences, Utility, and Demand: A Minnesota Symposium*, Harcourt Brace Jovanovich, 1971) seems to be one of the earliest publications with a proper axiomatic approach to preferences. Via an equivalence relation, two strict partial orders, and a completeness requirement, it defines a preference order as a total preorder. Note that, because cash flows add up, bundles are not necessary in financial contexts.

Summarily, it can therefore be assumed that there are two transitive relations which describe the typical finance models outlined above. First, an objective indifference relation, i.e. an equivalence relation, \sim' , can be used to describe which goods can be exchanged for one another in this market. Any market participant can be indifferent towards two goods $\alpha \sim' \beta$, because she or he can swap one for the other at any time. Note that I dropped any linearity assumption for greater generality. Second, there separately exists an objective strict preference order, \prec'' , in accordance to which all market participants act.

DEFINITION 22. *A financial market is given by a set of financial goods or cash flows, \mathcal{M} , together with an objective indifference relation, \sim' , on \mathcal{M} , which is an equivalence relation, and an objective strict preference order, \prec'' , on \mathcal{M} , which is a strict partial order.*

Two types of trades are possible for an agent in this market:

1. Exchange a cash flow for one which is considered equivalent to the first one by the market, i.e. receive α for β if $\alpha \sim' \beta$. Therefore, \sim' is the ‘fair exchange’ relation.
2. Trade a cash flow in for one which is considered strictly worse than the first one by the market, i.e. receive α for β if $\alpha \prec'' \beta$. Thus, \prec'' can be viewed as the ‘down-trade’ relation.

The following definition for the transitive closure of the two relations makes sense to me. Note that any transitive closure of a reflexive relation naturally is reflexive, too.

DEFINITION 23. *The preorder $\preceq_{\sim' \cup \prec''}$ generated by the fair exchange relation (objective indifference relation) and by the down-trade relation (objective strict preference order) is called the trade relation or the attainable trades.*

An important question in finance now is under which conditions a given market is rational in the sense that there exist no so-called arbitrage opportunities. These are trades, or chains of trades, where a cash flow can be swapped for an objectively better one at no additional cost. Such trades are also called a ‘free lunch’ (cf. [3]), or it is said that someone made ‘money out of nothing’. Simply put, these deals are ‘too good to be true’.

DEFINITION 24. *An arbitrage opportunity is a finite chain of either fair exchanges, down-trades, or both, which summarily exchanges a good $\alpha \in \mathcal{M}$ for an objectively better one, $\beta \in \mathcal{M}$, i.e. $\alpha \prec'' \beta$. The condition of ‘no-arbitrage’ is the absence of arbitrage opportunities in the market $(\mathcal{M}, \sim', \prec'')$.*

Without being able to here go into any details, in financial economics or mathematical finance, the so-called Fundamental Theorem of Asset Pricing (FTAP; e.g. [9], [4], and [3] in chronological order) establishes the equivalence of the absence of arbitrage and the existence of a risk-neutral pricing measure (RNPM; for stochastic one-period models) or an equivalent martingale measure (EMM; for multi-period or time-continuous models). By means of the risk-neutral pricing formula (a discounted expectation), the RNPM or the EMM implies a linear price functional for all cash flows which is in line, or consistent, with the existing prices and with the existing objective strict preference order (more preferable means more expensive). It therefore completes the market. Completeness here means that all cash flows are replicable by trading strategies and, thus, have a price. Having a price, however, means that any two cash flows can be compared with one another.

I think it is remarkable that, while later articles on FTAP are mostly concerned with the stochastic aspects, earlier publications (cf. [9], [4]) also show some focus on FTAP as an extension theorem of the pre-existing linear price functional. As such, also in its modern versions, FTAP is an extension or completion theorem and the *a priori* given linear trade relation and the strict preference relation can be completed in a consistent – which here means arbitrage-free – manner if and only if the original market is arbitrage-free. Moreover, FTAP also states that the RNPM or the EMM, and therefore the new price functional, is uniquely determined if and only if the *a priori* given market is complete.

THEOREM 3. *For a financial market $(\mathcal{M}, \sim', \prec'')$, the following are equivalent.*

1. *It is free of arbitrage opportunities.*
2. *Its fair exchanges, \sim' , and its down-trades, \prec'' , are chain-consistent with one another.*
3. *Its attainable trades consistently extend / are chain-consistent with its fair exchanges, \sim' , and its down-trades, \prec'' .*
4. *There exists a complete (math.: total) preference order on \mathcal{M} which consistently extends the objective indifference relation, \sim' , and the objective strict preference order, \prec'' .*

Under no-arbitrage, the attainable trades are complete (math.: total) if and only if the consistent completion is unique.

Proof. Def. 13 implies that the chain-consistency of $\preceq_1 = \sim'$ and $\preceq_2 = \prec''$ is equivalent to no-arbitrage (note that here $\prec_1 = \sim_2 = \emptyset$). This is because on the one hand, an arbitrage trade chain establishes a closed chain as in Def. 13, where the at least one $(\gamma_i, \gamma_{i+1}) \in \prec_1 \cup \prec_2$ is warranted by $\alpha \prec'' \beta$. Vice versa, if a chain as in Def. 13 exists, then the at least one $(\gamma_i, \gamma_{i+1}) \in \prec_1 \cup \prec_2$ in fact implies at least one $(\gamma_i, \gamma_{i+1}) \in \prec''$, so I can set $\alpha = \gamma_i$ and $\beta = \gamma_{i+1}$. Theorem 3 now is a direct consequence of Corollary 6. \square

For a distinction of notions, I mention that in [6], a ‘complete’ preference order is called ‘rational’. In mathematics, however, the descriptors ‘total’ (Def. 1) or ‘strongly connected’ are used if all elements of \mathcal{M} can be compared using the considered preorder. The notion ‘rational’ I would rather use interchangeably to ‘arbitrage-free’.

Like FTAP, Theorem 3 is an extension theorem. It equates no-arbitrage, which is the consistency of two pre-existing preference orders on a part of the cash flows, to the existence of a preference order (FTAP: price functional/RNPM/EMM) for all cash flows which is consistent with the two pre-existing preference orders. Moreover, if this extension of the preference orders is uniquely determined, the market has a completeness property in the sense that all cash flows were comparable *a priori* and could be traded for one another in at least one direction. Hence, the parallels of FTAP to Theorem 3 are striking, and while my result stays clear of any linearity assumptions or stochastic models and therefore does not directly imply FTAP in its modern form, I consider Theorem 3 a broad generalization of it.

9 Application in microeconomic theory

In this section, I assume that the preferences of two entities, e.g. of two individuals, named 1 and 2, are given by the additive positively homogeneous preorders, i.e. vector preorders, \preceq_1 and \preceq_2 on the real linear space \mathcal{M} with the nullvector o . For instance, one could think of preferences relating to appropriately defined deterministic or stochastic cash flows. Additivity, positive homogeneity, and reflexivity at first, and superficially, appear to be reasonable properties of preferences in a financial context.

If individual 1 holds good $\alpha \in \mathcal{M}$ and 2 holds $\beta \in \mathcal{M}$, then $\alpha + \beta \in \mathcal{M}$ is the total amount in goods (e.g. the total cash flow) that is allocated to 1 and 2. If for $\alpha', \beta' \in \mathcal{M}$ with $\alpha' + \beta' = \alpha + \beta$ it holds that $\alpha \preceq_1 \alpha'$ and $\beta \preceq_2 \beta'$, and at least one \preceq_i for $i = 1, 2$ can be replaced with its strict version, \prec_i , then (α', β') is called a Pareto improvement of the allocation (α, β) . As is well-known, a Pareto improvement therefore improves the situation of at least one individual while not worsening the situation of the other.

THEOREM 4. *Let \preceq_1 and \preceq_2 be two additive positively homogeneous preorders (vector preorders). There exists a Pareto improvement for any allocation between two entities with these preferences if and only if \preceq_1 and \preceq_2 are not chain-consistent with one another.*

By Theorem 2, under the conditions of Theorem 4, chain-consistency is equivalent to the existence of a total additive positively homogeneous preorder which consistently extends both, \preceq_1 and \preceq_2 . So, expressed differently, *if two individuals with additive positively*

homogeneous preferences are not both ‘cut from the same cloth’, there exists no Pareto optimal allocation, and, in this sense, no equilibrium. Obviously, an analog statement follows for an individual with preferences \preceq_1 vis-à-vis an entire market, if the attainable trades of this market as in Def. 23 of Sec. 8 are given by \preceq_2 . In a loose sense, my result adds to the row of impossibility problems concerning preference orders in microeconomics, such as Arrow’s Impossibility Theorem or the Condorcet Paradox (cf. [6]).

Proof of Theo. 4. Generally, for the following, note that for any additive preorder \preceq on a linear space \mathcal{M} , $\alpha \preceq \beta$ implies $\alpha - \beta \preceq o$ and $-\beta \preceq -\alpha$, since, by additivity and reflexivity, I can consecutively add $-\beta \preceq -\beta$ and then $-\alpha \preceq -\alpha$ to $\alpha \preceq \beta$. Now, if (α', β') is a Pareto improvement of some $(\alpha, \beta) \in \mathcal{M} \times \mathcal{M}$, then $\alpha' = \alpha + \delta$ and $\beta' = \beta - \delta$ for some $\delta \in \mathcal{M}$, since the total allocation does not change. With the rules explained earlier and since $\alpha \preceq_1 \alpha + \delta$ and $\beta \preceq_2 \beta - \delta$, a Pareto improvement therefore is equivalent to the existence of some $\delta \in \mathcal{M}$ with

$$o \preceq_1 \delta \preceq_2 o, \quad (49)$$

such that at least one \preceq_i for $i = 1, 2$ in (49) can be replaced with its strict version, \prec_i .
 \Rightarrow : By Def. 13, if the preferences are consistent, then a closed chain of length two as just explained in (49) is impossible. Thus, a Pareto improvement implies inconsistency.
 \Leftarrow : Inconsistency implies for some $n \in \mathbb{N} \setminus \{0, 1\}$ that $\gamma_i \preceq_1 \gamma_{i+1}$ for all $i \in I$ and that $\gamma_i \preceq_2 \gamma_{i+1}$ for all $i \in J$, where I and J are disjoint and nonempty (cf. Lemma 9), $I \cup J = \{1, \dots, n-1\}$, and $\gamma_1 = \gamma_n$. Additivity and reflexivity of \preceq_1 and \preceq_2 now imply

$$o \preceq_1 \sum_{i \in I} (\gamma_{i+1} - \gamma_i), \quad (50)$$

as well as

$$\sum_{i \in J} (\gamma_i - \gamma_{i+1}) \preceq_2 o. \quad (51)$$

Because of the requirements of Def. 13 and by the second statement of Lemma 3, (50), (51), or both, must be strict. However,

$$\begin{aligned} \sum_{i \in J} (\gamma_i - \gamma_{i+1}) &= \sum_{i=1}^{n-1} (\gamma_i - \gamma_{i+1}) - \sum_{i \in I} (\gamma_i - \gamma_{i+1}) \\ &= \gamma_1 - \gamma_n + \sum_{i \in I} (\gamma_{i+1} - \gamma_i). \\ &= \sum_{i \in I} (\gamma_{i+1} - \gamma_i), \end{aligned} \quad (52)$$

and, with $\delta := \sum_{i \in I} (\gamma_{i+1} - \gamma_i)$, one obtains a chain as in (49) from (50) and (51). \square

In financial economics, or mathematical finance, creating a proper financial advantage out of nothing is called a ‘free lunch’ or an ‘arbitrage’ (see also Sec. 8). Modern financial mathematics is based on the principle of ‘no-arbitrage’, the exclusion of arbitrage for the reason that such opportunities quickly disappear in real markets. In finance, it matters

what exactly ‘proper’ means, because there it needs to be an objective notion, while in the situation at hand here, it does not necessarily, since I can consider the preferences of two individuals, and all that matters is what the individuals prefer over the zero, o , and not, for instance, what a market summararily prefers.

Reconsidering the property of additivity, in a financial context, this may not be the wisest of properties for individual preferences. In risk management or insurance, diversification, which can stem from stochastic independence, has always been a very important aspect. If $\alpha \preceq_i \beta$ and $\alpha' \preceq_i \beta'$, but α and α' enabled diversification by independence, while β and β' do not (and maybe do quite the opposite when lumped together), preferring $\beta + \beta'$ over $\alpha + \alpha'$ could be a bad idea. Theorem 4 provides a further reason why, generally, additive preferences are problematic and, so I presume, unrealistic.

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